

What is this “Viterbi Decoding”

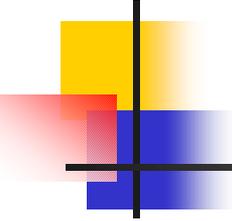
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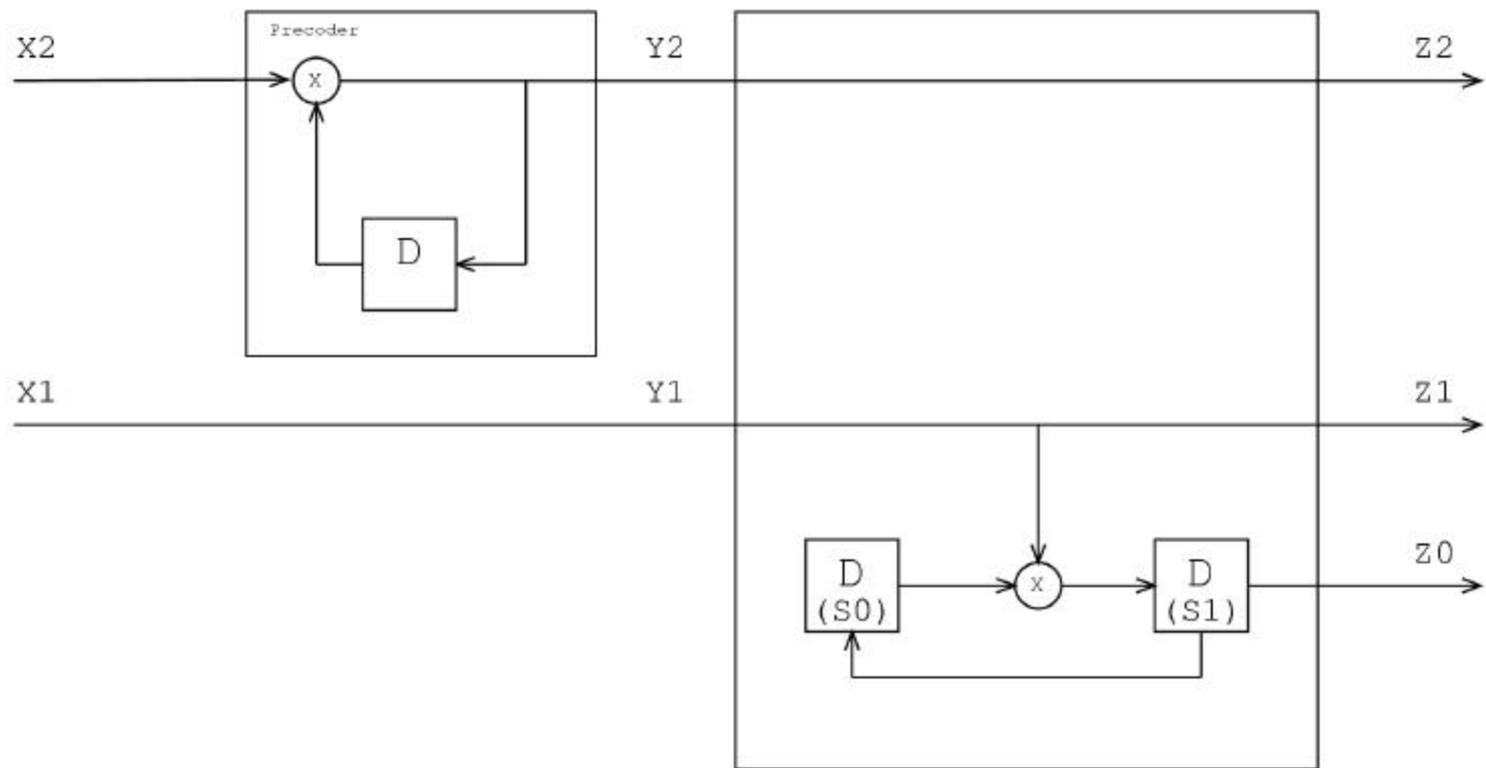
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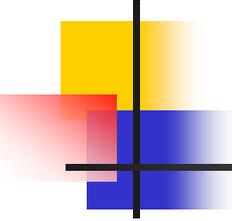
State Machine/Trellis Encoding?

- You can also represent the convolution encoder as a finite state machine
- The possible state evolution of the convolution encoder traces out a “trellis” like structure in state-time space.

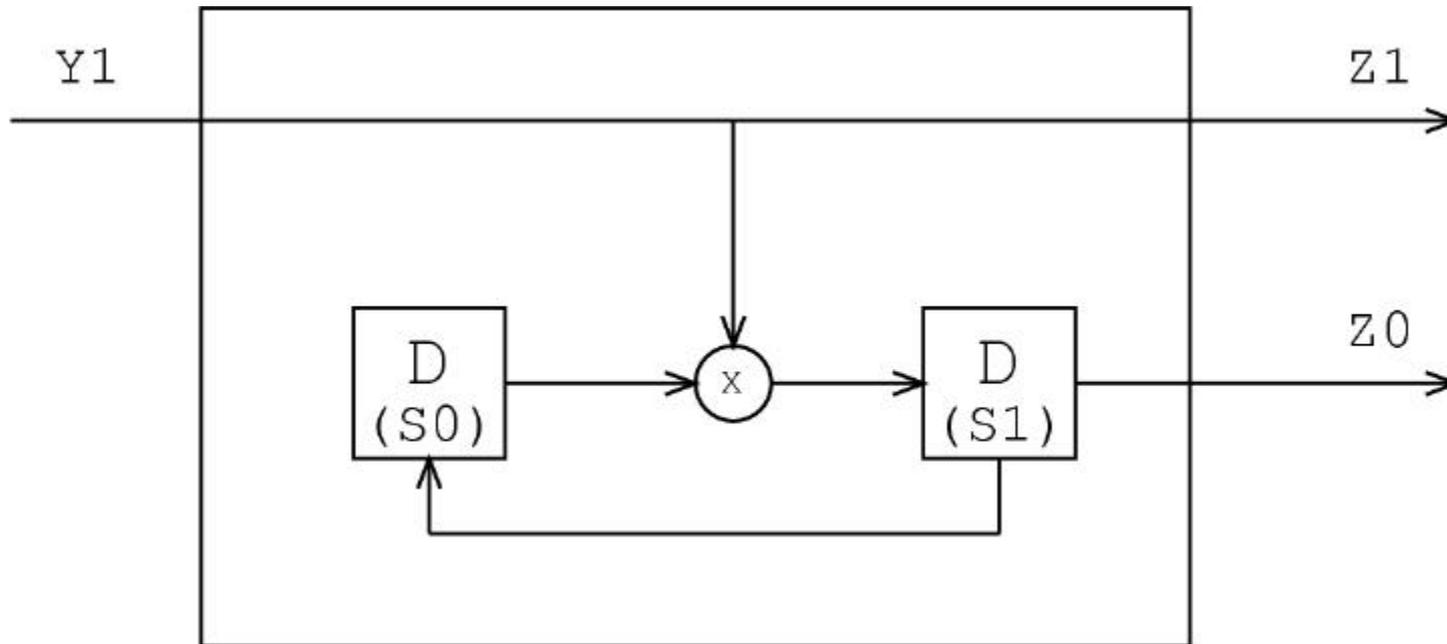
Convolutional Encoder



Source: ASTC Standard A53, revision B



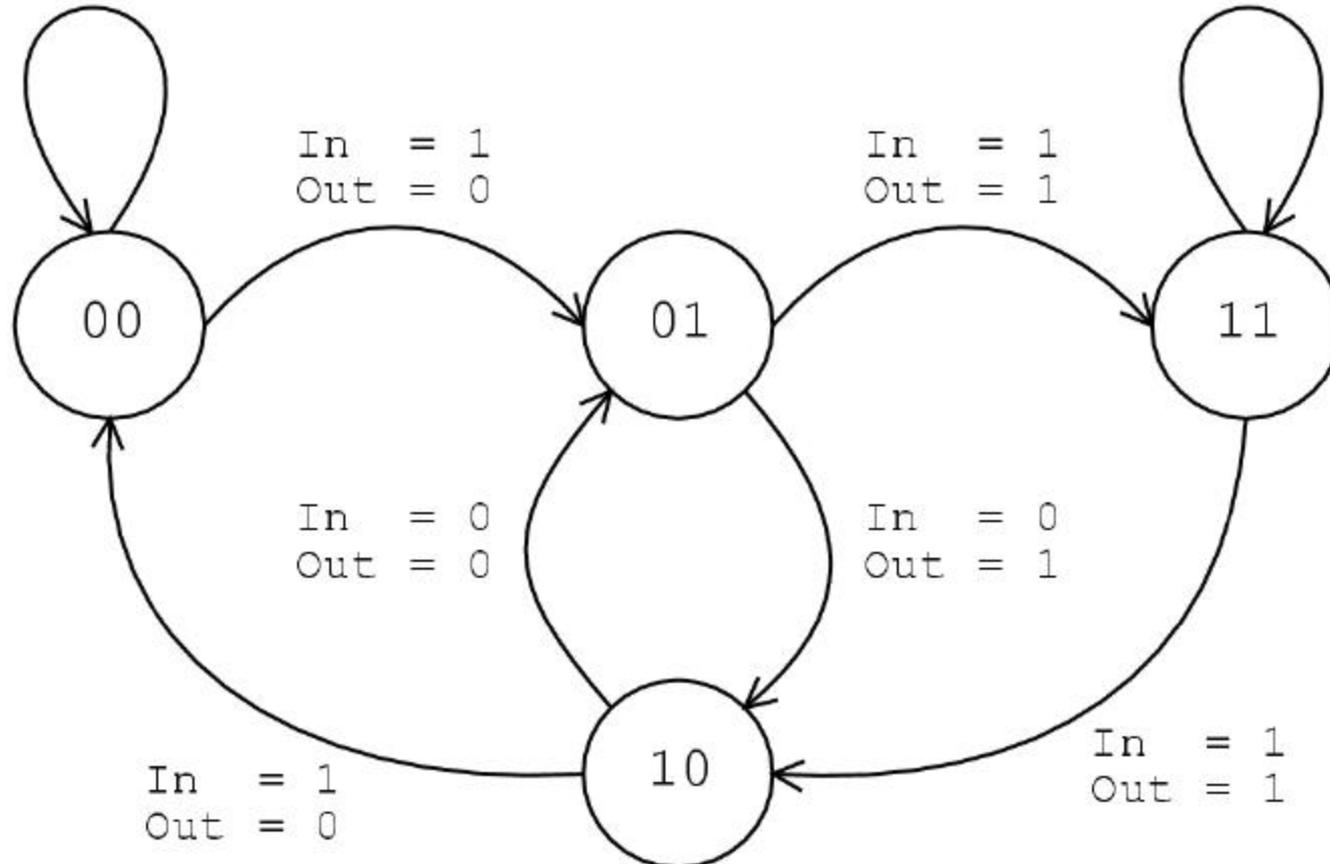
Interesting Part



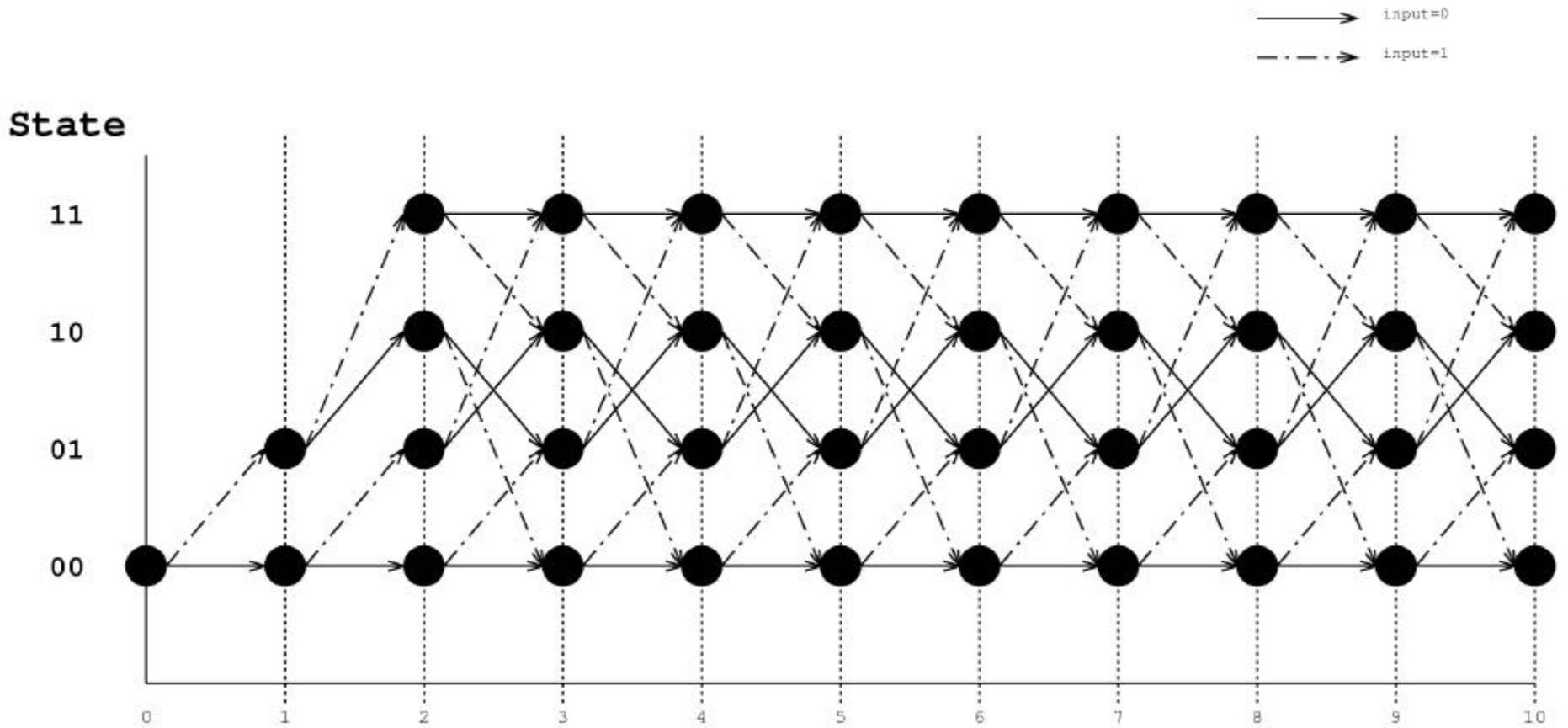
State Transition Diagram

In = 0
Out = 0

In = 0
Out = 1



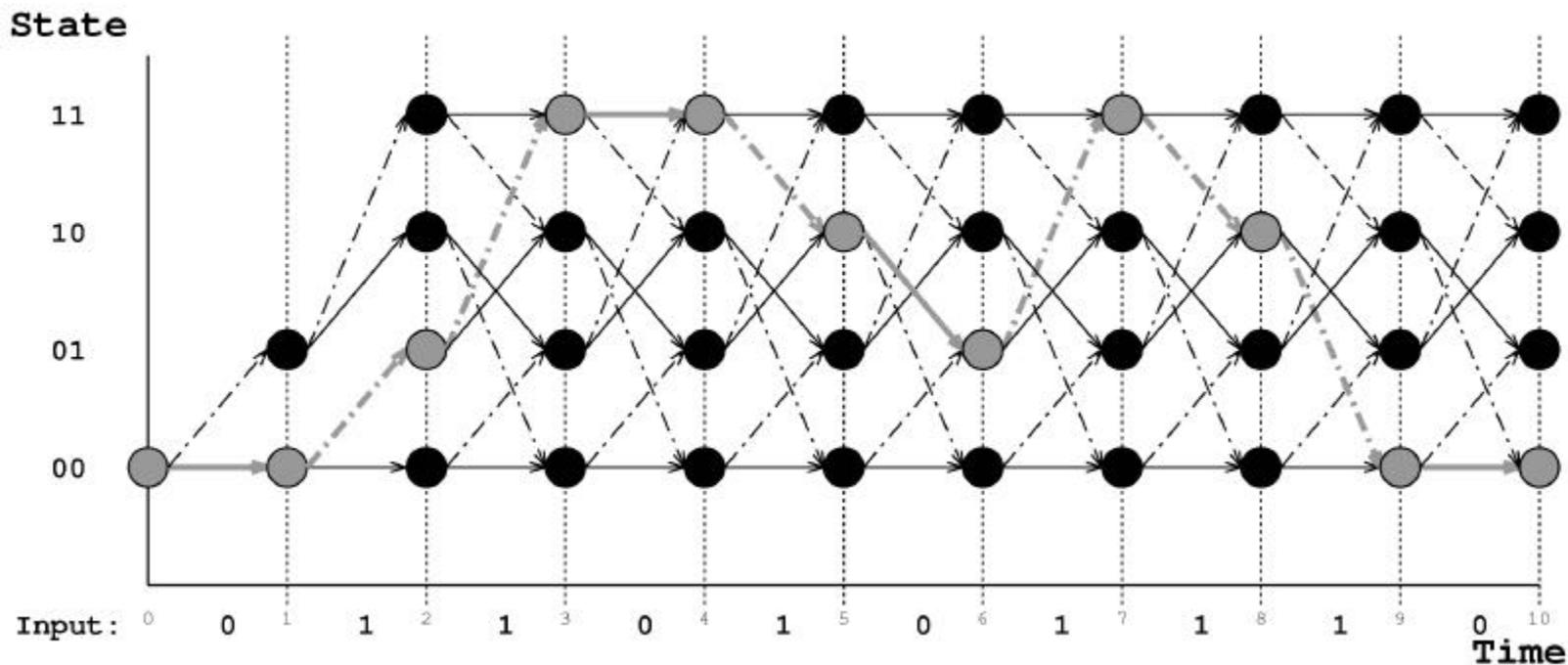
Trellis Diagram



Note: it is assumed that you start in state 00

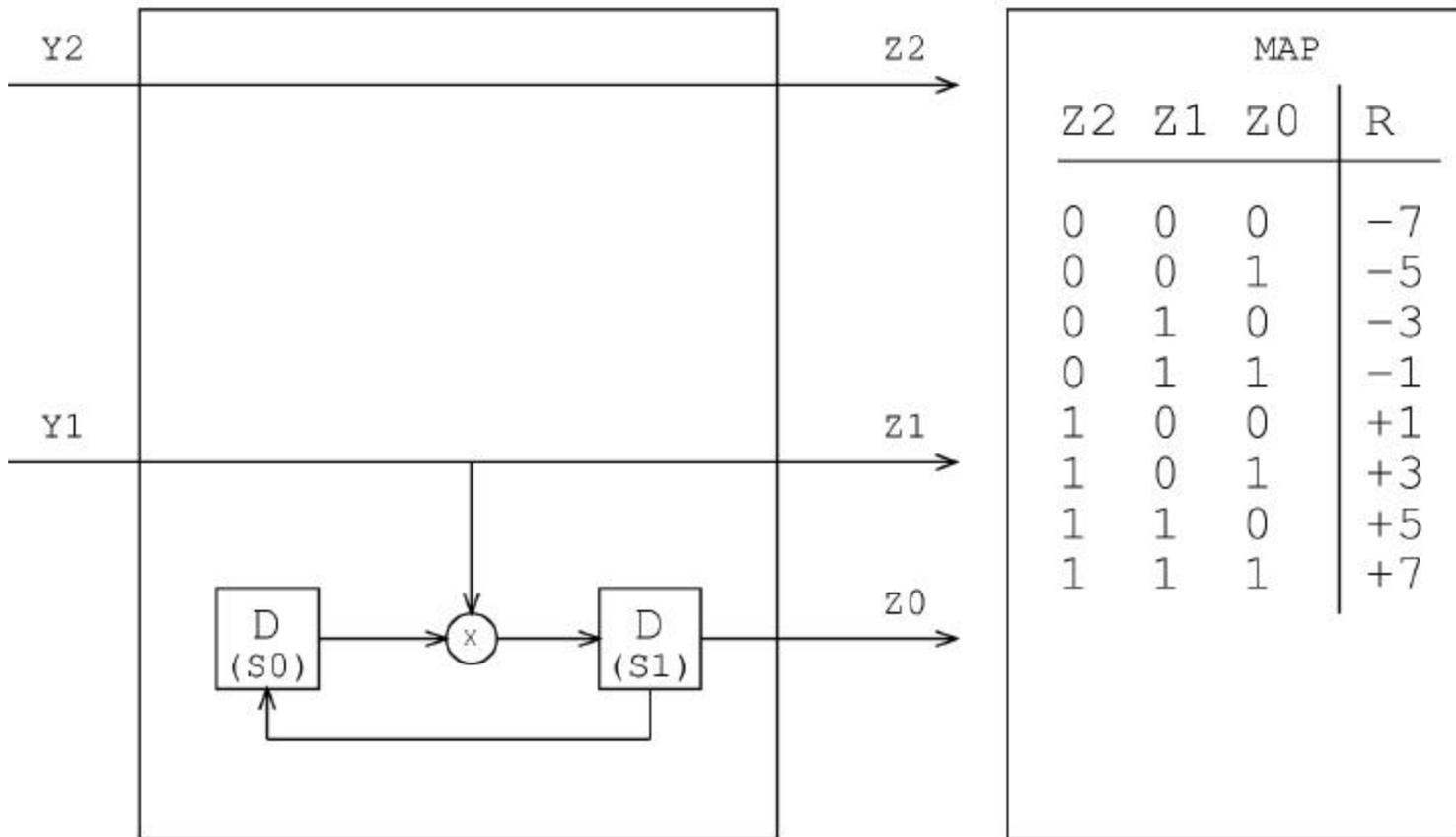
Trellis Paths

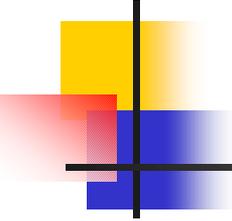
- Each set of inputs traces a unique path through the trellis.



Input:
01101
01110

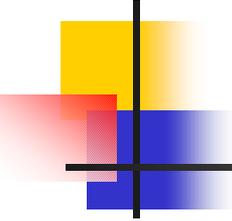
Symbol Mapping





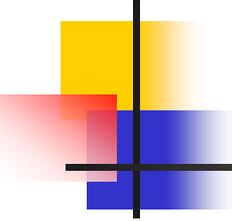
Trellis Decoder

- Use Viterbi Algorithm to undo Convolution Coding (not signal \rightarrow symbol unmapping)
- Two flavors of determining unmapping
 - Hard: make a decision as to which symbol the received signal most closely matches.
 - Soft: Assign weights to all symbols based on their respective likelihood given received signal
- Viterbi implementation we present assumes a hard decision model.



Viterbi Decoding (Overview)

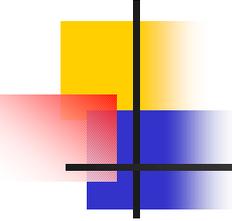
- The Viterbi algorithm:
 - Given a sequence of received symbols, (that were produced by a convolution encoder, sent over a channel)
 - Determine what the input to the convolution encoder was
- It does this by determining the most likely path through the trellis



Viterbi Decoding (Main Idea)

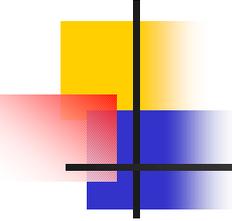
- Dynamic Programming
 - Keep a table $\mathbf{c}[s,t]$ that records the number of errors* that would have been accumulated if the encoder was in state s at time t .
 - Also keep a table $\mathbf{p}[s,t]$ which records the state that the encoder would have been in at time $t-1$ if it were in state s at time t .

* Typically calculated using either Hamming or Euclidian distance



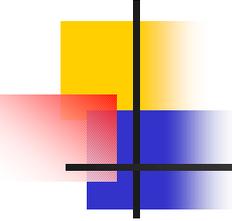
Filling out the tables

- At time t , we receive the symbol R_t .
- For each state s ,
 - Let q_0 and q_1 be the two possible previous states of s at time $t-1$.
 - Let e_0 be the error between R_t and $q_0 \rightarrow s$
 - Let e_1 be the error between R_t and $q_1 \rightarrow s$
 - $c[s,t] \leftarrow \min(\mathbf{c}[q_0,t-1]+e_0, \mathbf{c}[q_1,t-1]+e_1)$
 - Update $\mathbf{p}[s,t]$ appropriately with q_0 or q_1



Traceback

- When the algorithm has examined T input symbols, it looks for the minimum entry among all states in $\mathbf{c}[s, T]$.
- Then the algorithm traces back through the trellis using the entries of $\mathbf{p}[s, t]$.



Example

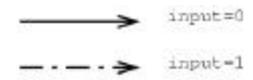
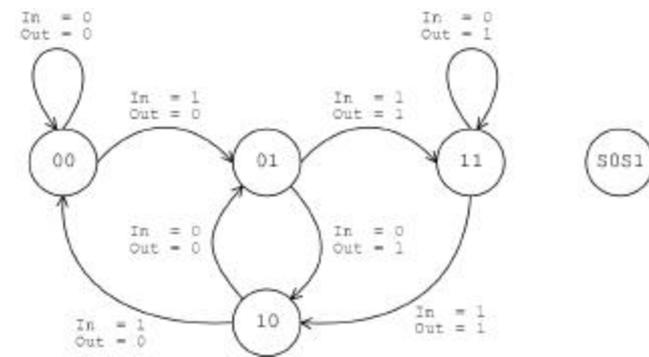
- Input to Encoder:

- 0110101110

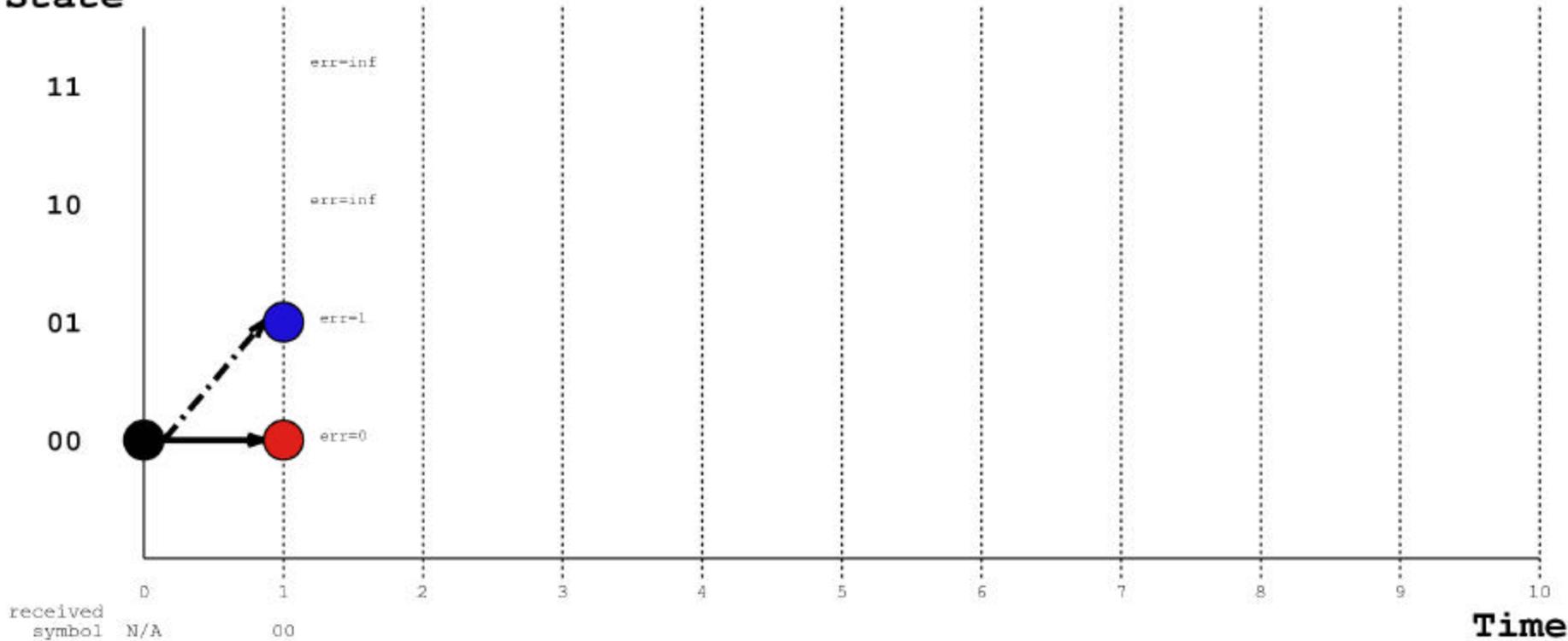
- Output of Encoder:

- 00 10 11 01 11 00 11 11 10 00

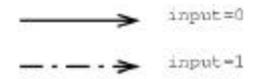
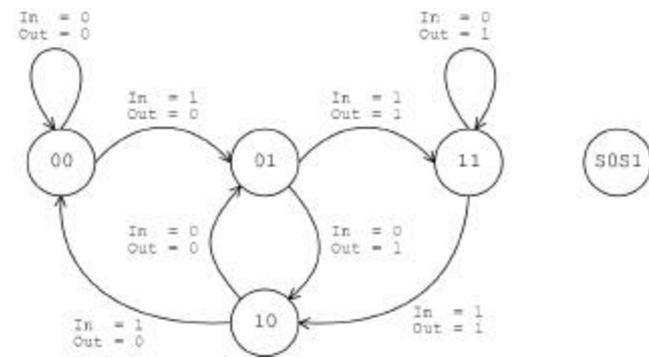
Decoder, time = 1



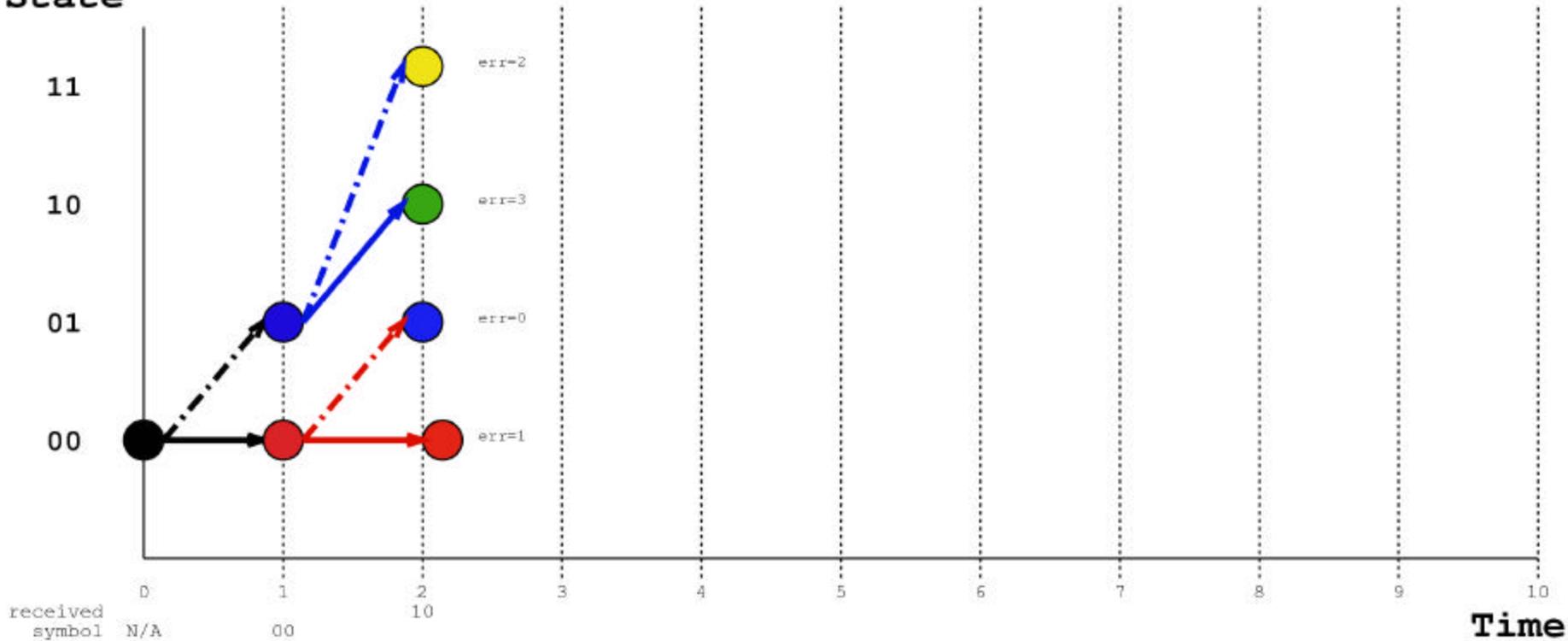
State



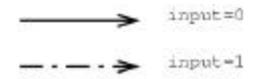
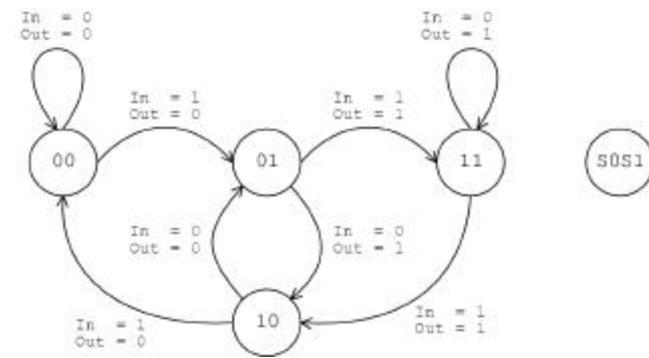
Decoder, time = 2



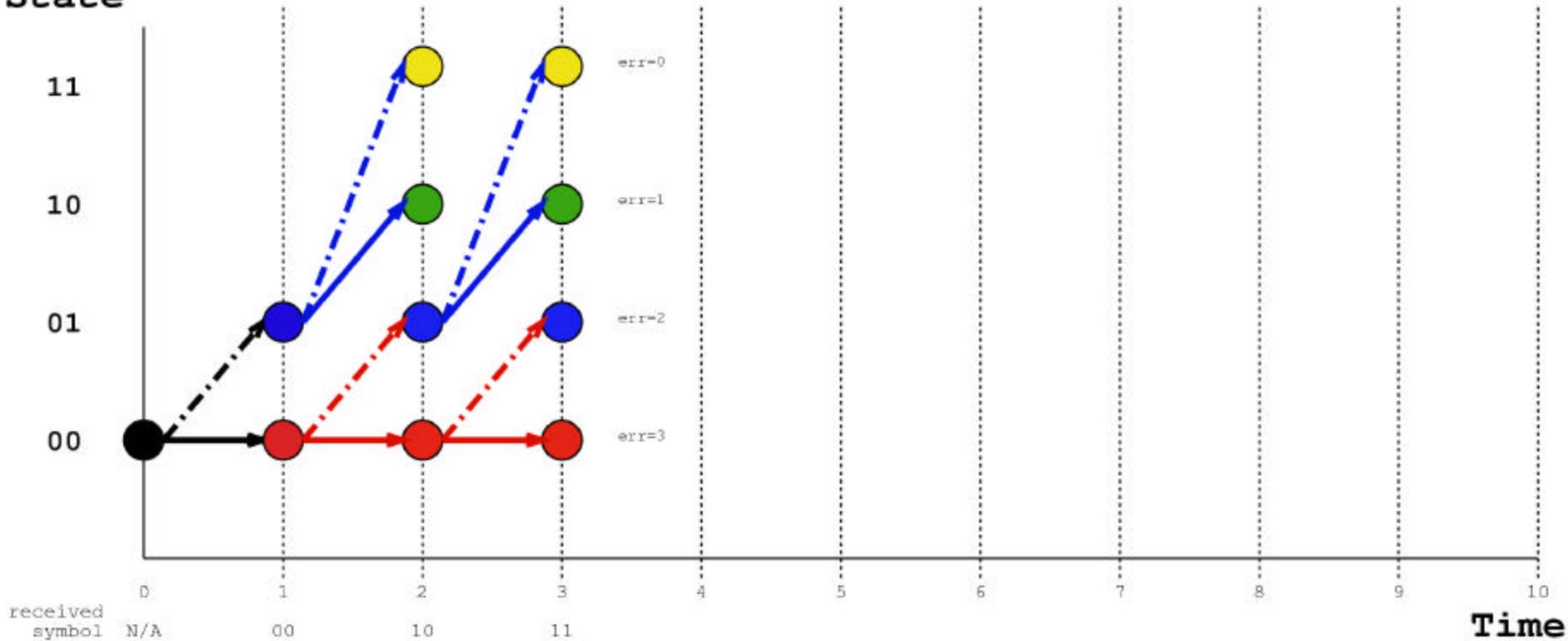
State



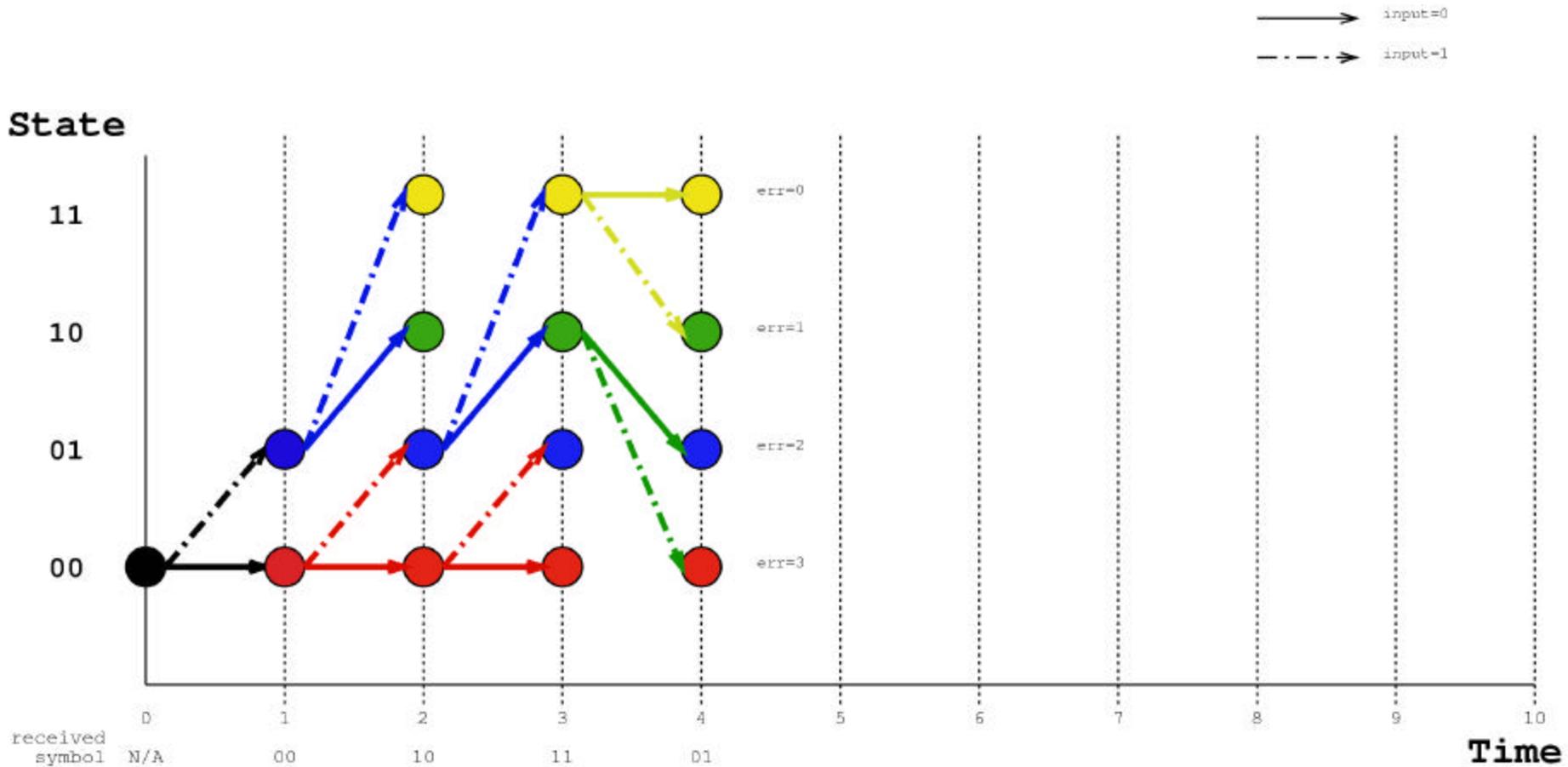
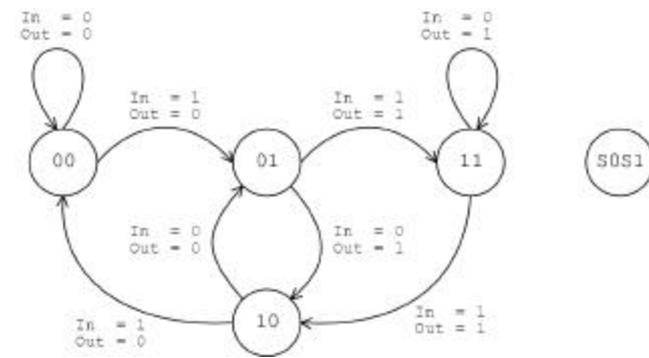
Decoder, time = 3



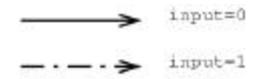
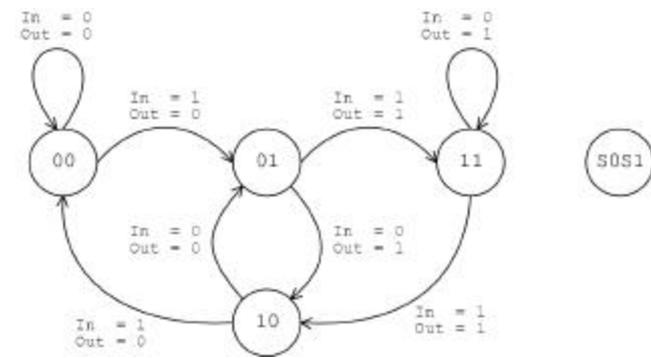
State



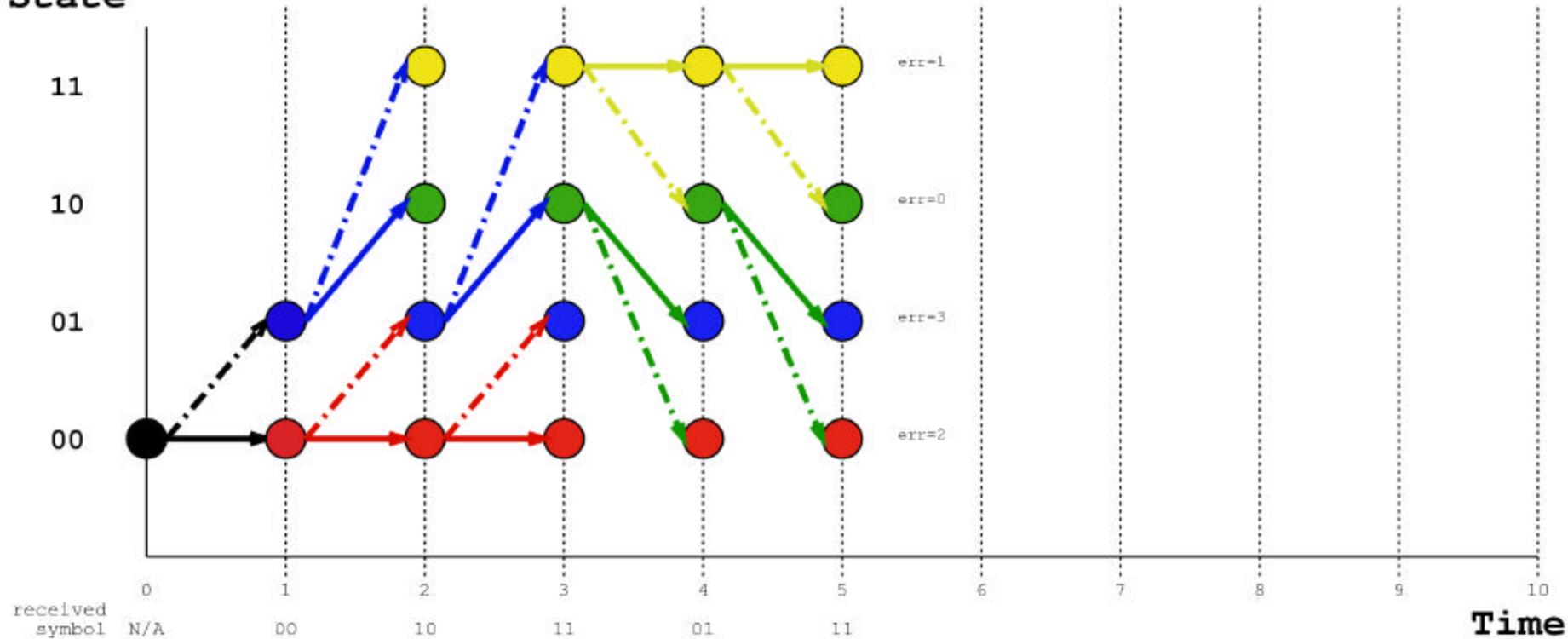
Decoder, time = 4



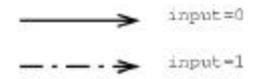
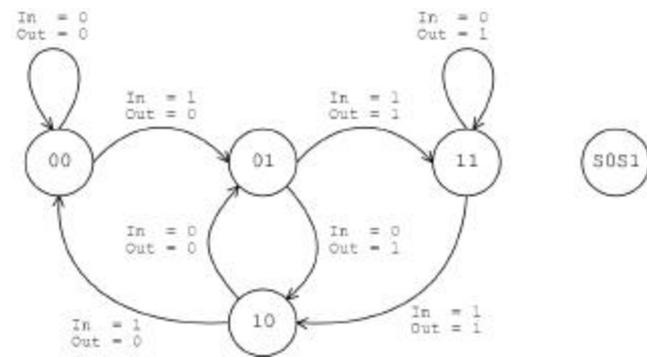
Decoder, time = 5



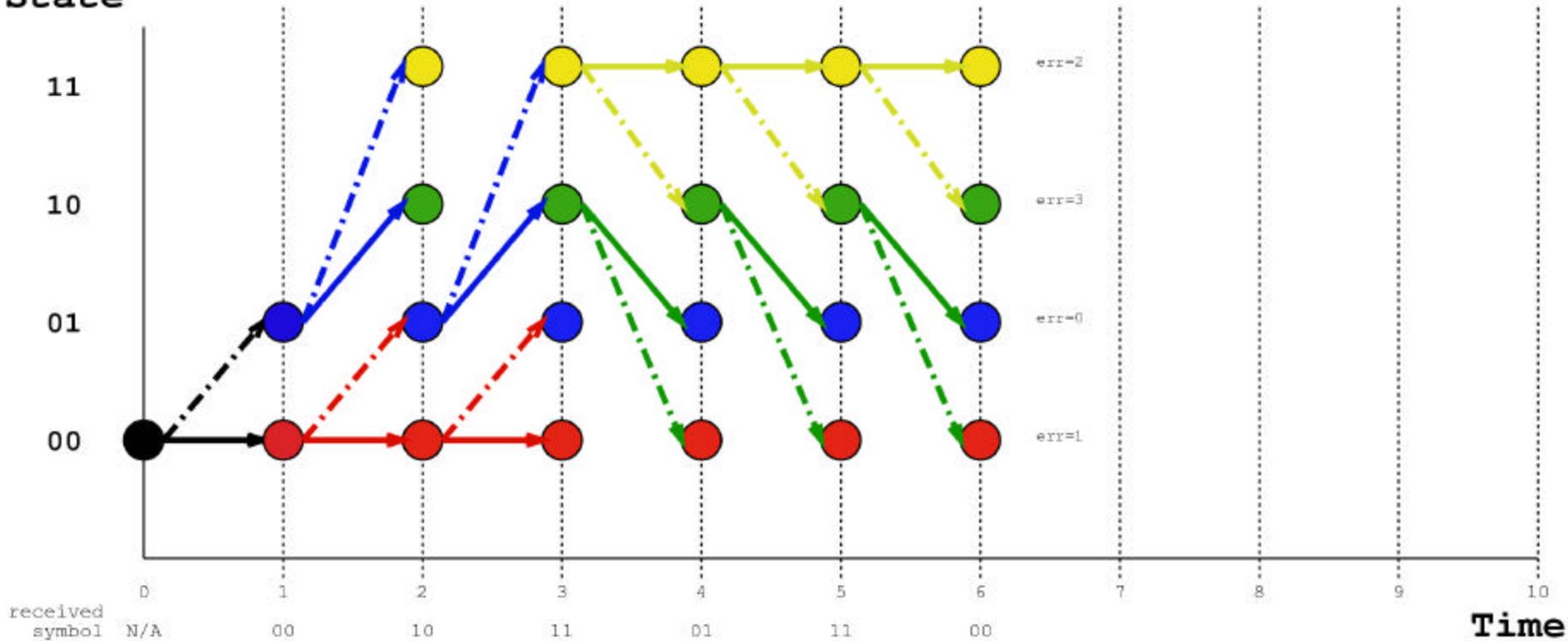
State



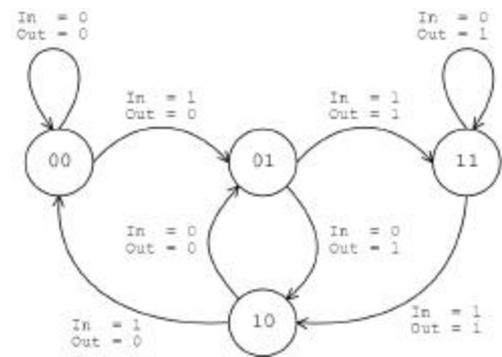
Decoder, time = 6



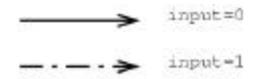
State



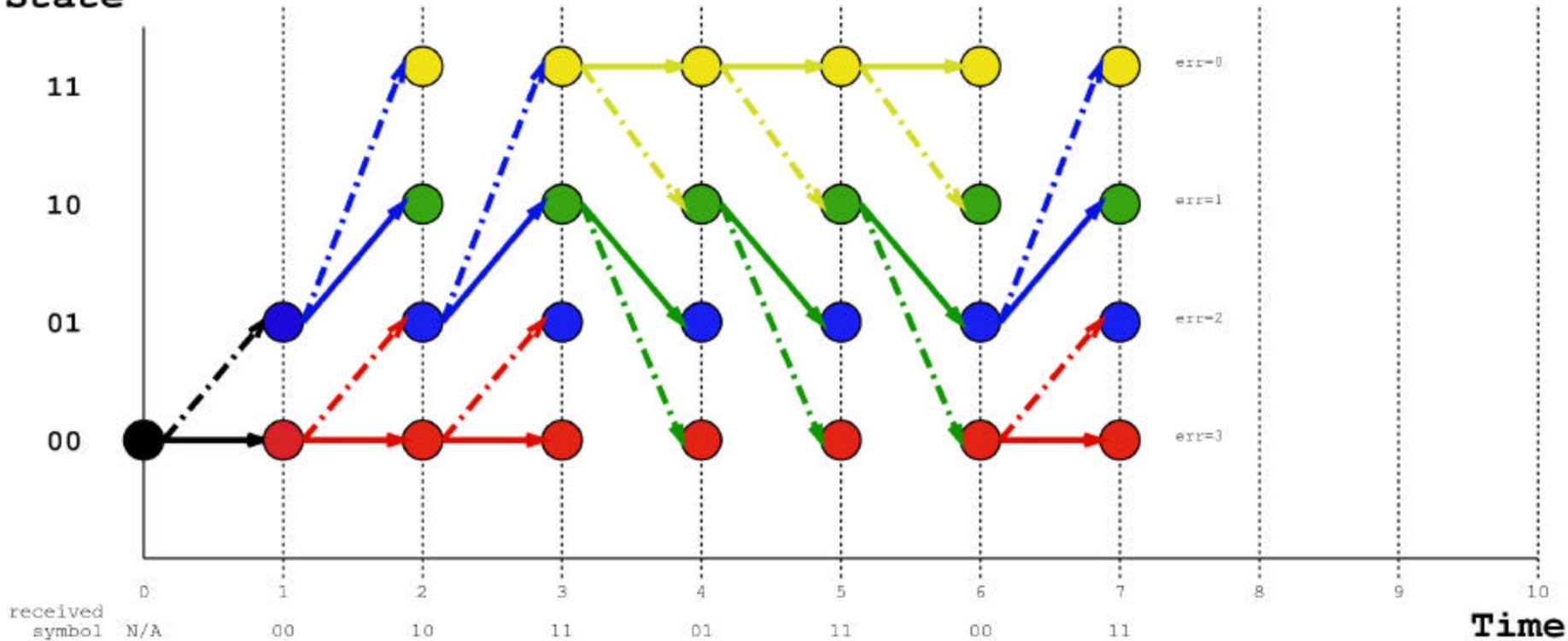
Decoder, time = 7



S0S1

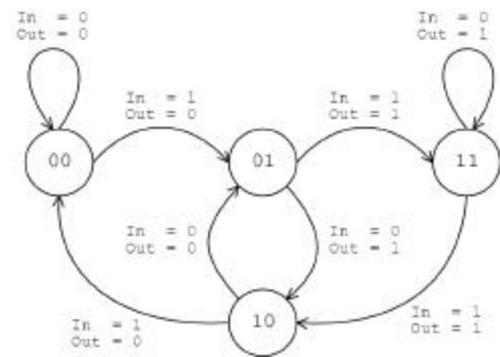


State

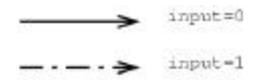


Time
21

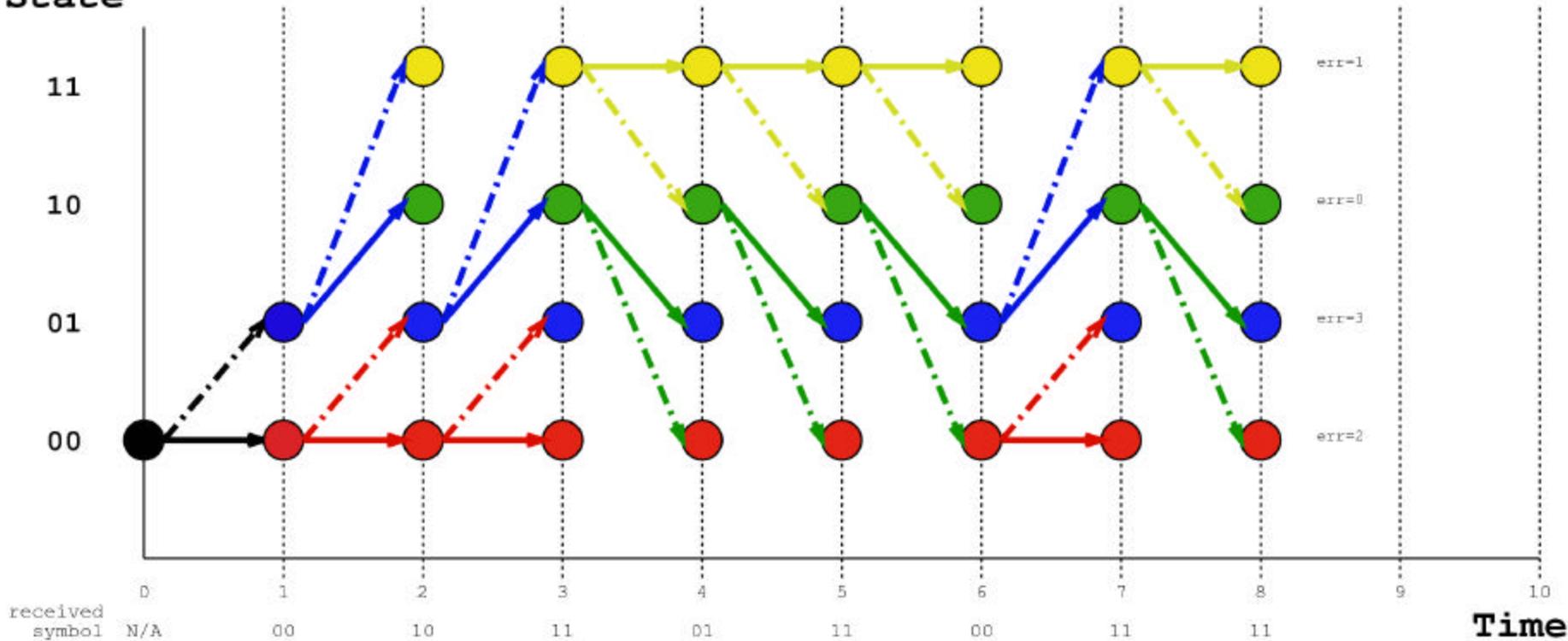
Decoder, time = 8



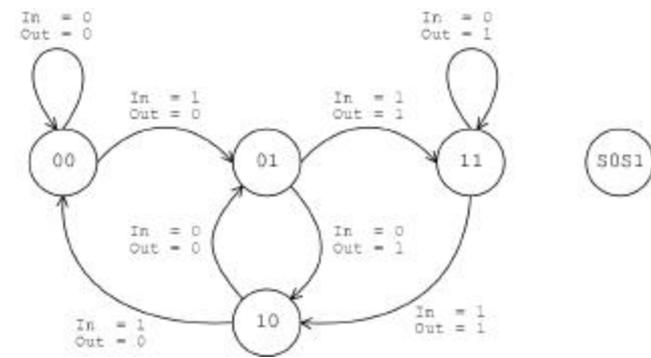
S0S1



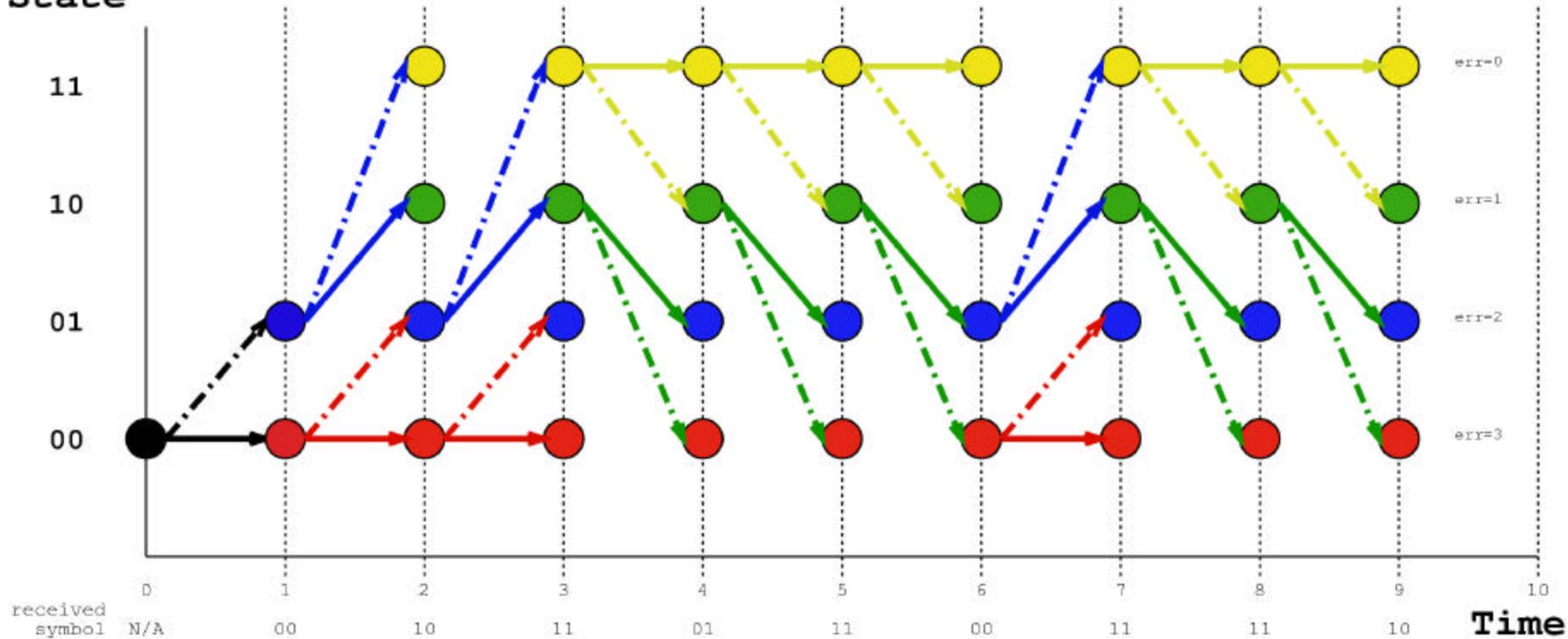
State



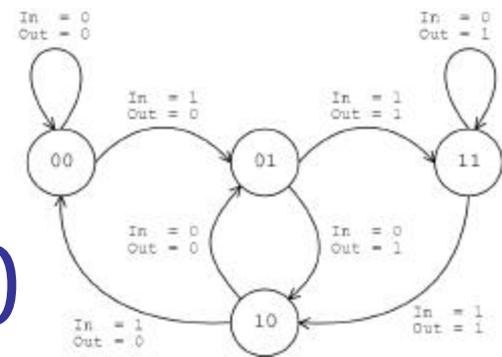
Decoder, time = 9



State



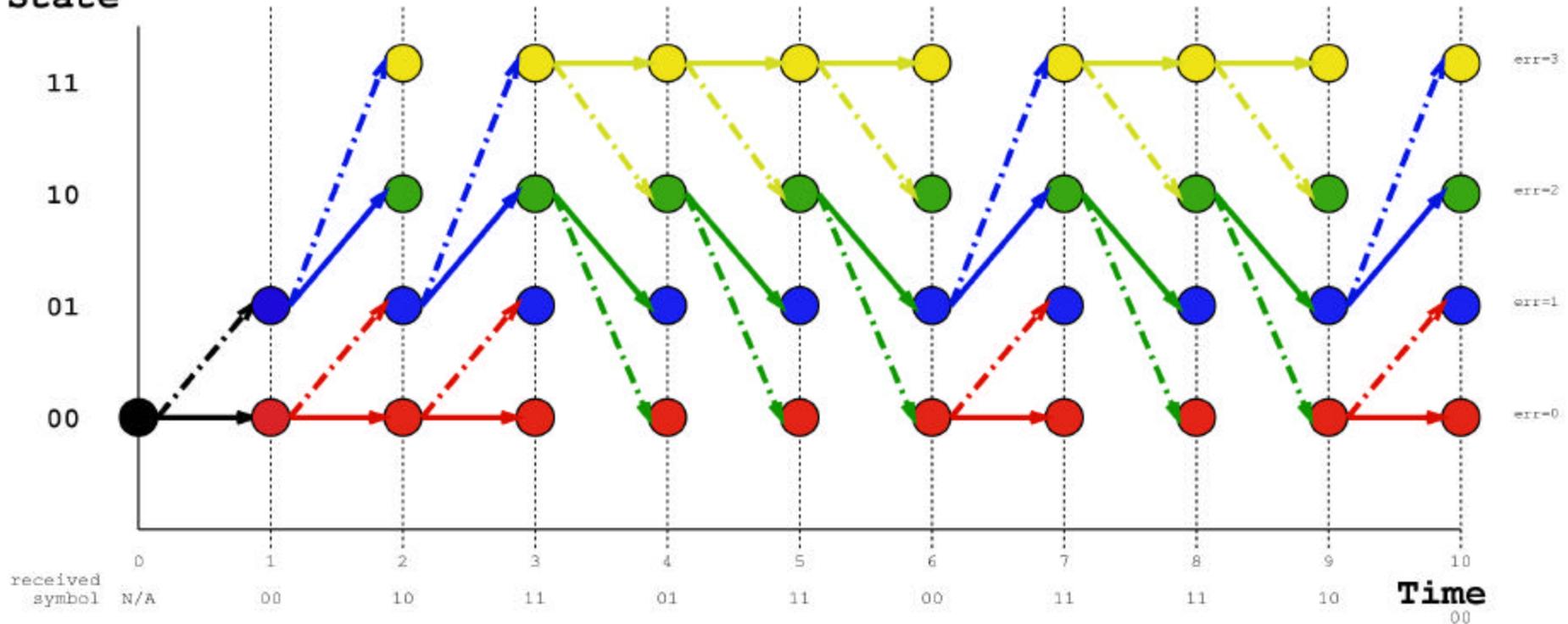
Decoder, time = 10



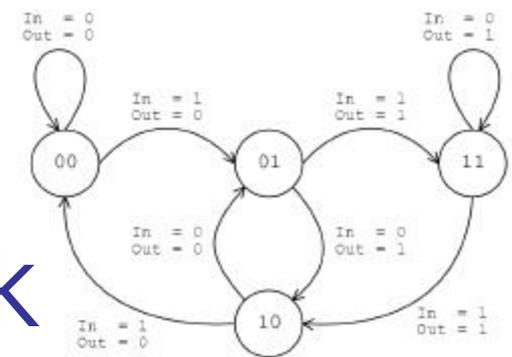
S0S1



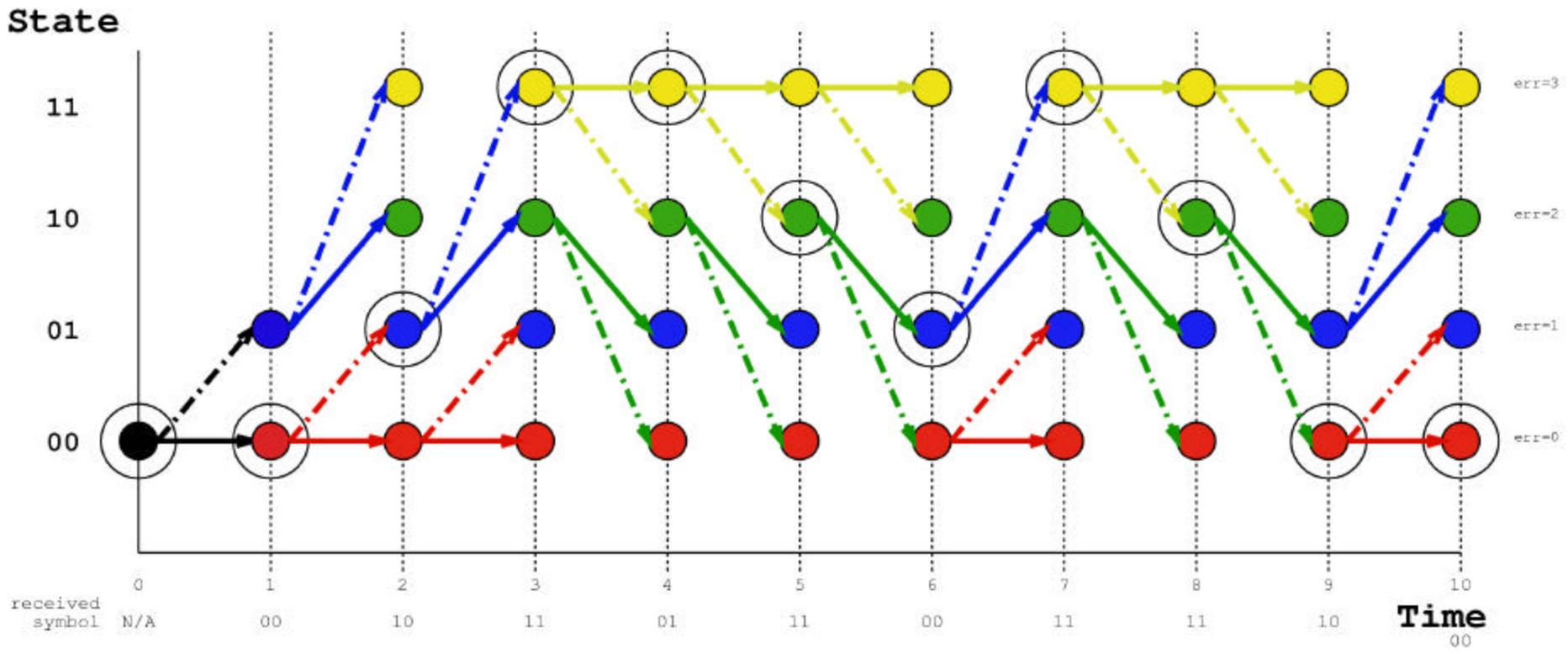
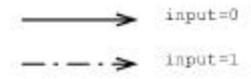
State



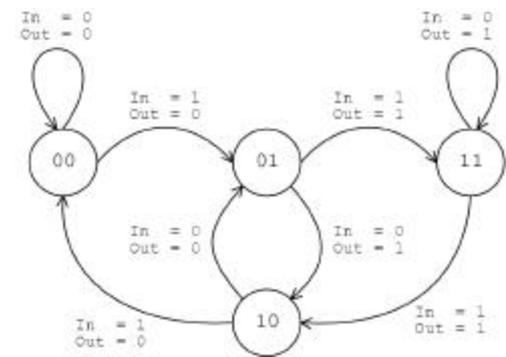
Decoder, Traceback



SOS1



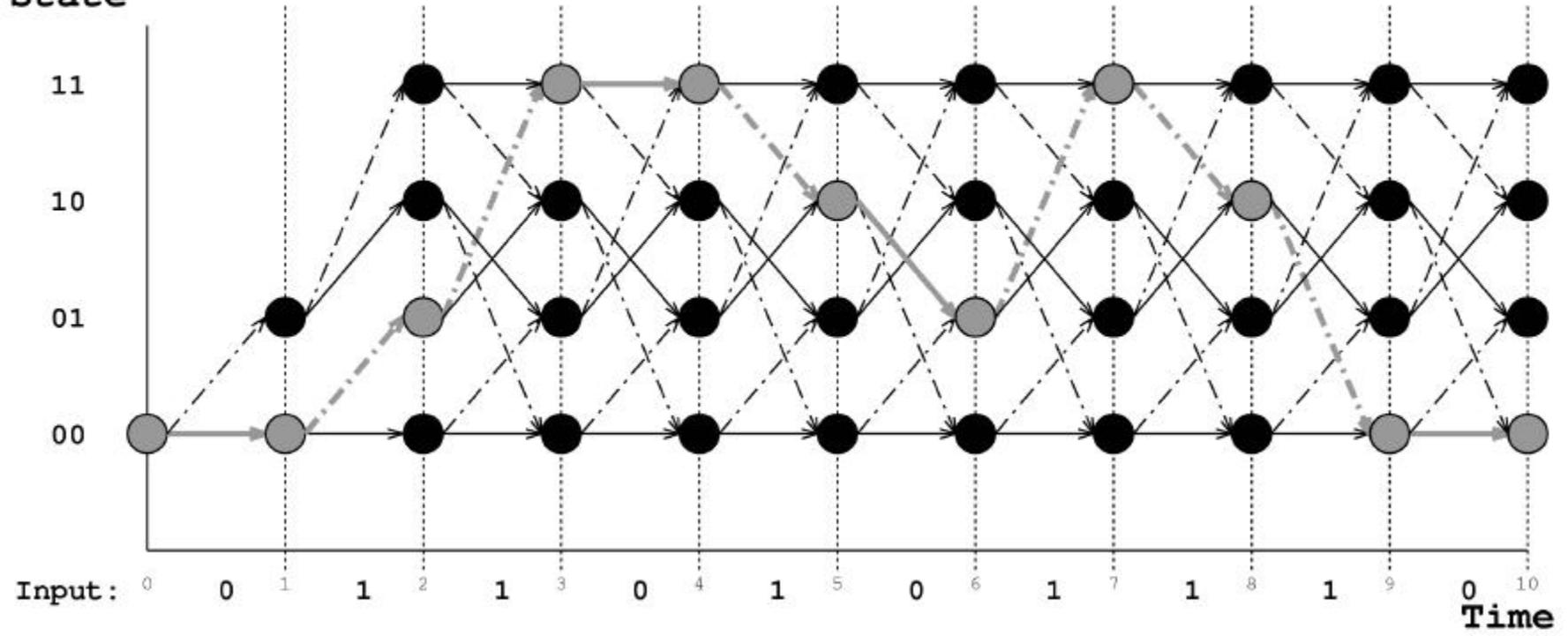
Original Path

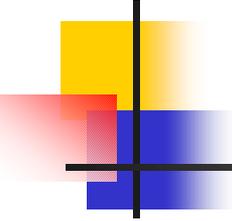


S0S1



State





Questions

- Is there a way to express the Viterbi algorithm in a fine grained stream graph?
- Variable rates for this decoder would mean `pop(2)` for 5K times, and then `push(5K)`. This is deterministic, if not constant rate.
- Perhaps 2 stage filters?